Lecture 8

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2. Little's Law

- > one of the most important results in queuing theory;
- ➤ it was first established and used by Philip. M. Morse and other researchers in the 1950's; in 1954, Morse published it, but was not able to prove it, so he challenged his readers to find a situation where it did not hold;
- ➤ **John D. C. Little** (1925 2024), Professor Emeritus at the MIT Sloan School of Management, proved it in 1961; later, in the 1990's and 2000's there were more developments and versions both in theory and in practice;
- ➤ Little's Law gives a simple relationship between the expected number of jobs, the expected response time, and the arrival rate; it is valid for any stationary queuing system.



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Proposition 2.1 (Little's Law).

In any stationary queuing system, the following holds:

$$E(X) = \lambda_A E(R). \tag{2.1}$$

Example 2.2.

A person walks into a bank at 10:00 a.m. He counts a total of 10 customers in the bank and assumes that this is the typical, average number. He also notices that, on average, new customers walk in every 2 minutes. When should he expect to finish his business and leave the bank?



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Solution.

The average number of customers in the bank, i.e. the expected number of jobs in the system, is

$$E(X) = 10.$$

On average, new customers walk in every 2 minutes, that is the mean interarrival time,

$$\mu_A = 2 \text{ minutes, so}$$
 $\lambda_A = 1/\mu_A = 1/2 / \text{ minute.}$

Then the amount of time he is expected to spend in the bank, i.e. the expected response time is, by Little's Law,

$$E(R) = \frac{1}{\lambda_A} E(X) = \mu_A E(X) = 20 \text{ minutes.}$$

Thus, he should expect to leave at 10:20.



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Remark 2.3.

1. Little's Law is universal, it applies to any stationary queuing system and even to the system's components, the queue and the servers.

Thus, we can immediately deduce the equations for the number of waiting jobs,

$$E(X_w) = \lambda_A E(W),$$

and for the number of jobs currently receiving service,

$$E(X_s) = \lambda_A E(S).$$

Note that the same arrival rate, λ_A , applies to the components, as for the entire queuing system.



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Remark (Cont).

2. Looking at the second equation above, E(S) is the expected or the mean service time, i.e. μ_S . So, we have

$$E(X_s) = \lambda_A \cdot \mu_S = \frac{\lambda_A}{\lambda_S} = r,$$

so we just obtained another important definition of utilization, which also justifies its name.

Definition 2.4.

Utilization r is the expected number of jobs receiving service at any given time.



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- ➤ Little's Law only relates expectations of the number of jobs and their response time;
- in the remaining sections of this chapter, we evaluate the entire distribution of X(t), which will help us compute various probabilities and expectations of interest;
- these quantities will describe and predict the performance of a queuing system.



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Definition 2.5.

The number of jobs in a queuing system, X(t), is called a **queuing process**.

- since X(t) is the number of jobs in the system, it is clearly a discrete-state stochastic process; the time set may be discrete or continuous and we will look at both cases;
- in general, a queuing process is not a counting process because jobs arrive and depart, therefore, their number may increase and decrease, whereas any counting process is nondecreasing;
- however, we will use counting processes to model arrivals and service of jobs;
- ➤ another aspect is the number of servers in a queuing system, one or more. Again, we will consider both situations (in the end, even considering the case where the number of servers goes to infinity).



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3. Bernoulli Single-Server Queuing Process

Definition 3.1.

A Bernoulli single-server queuing process (B1SQP) is a discrete-time queuing process with the following characteristics:

- one server:
- unlimited capacity;
- arrivals occur according to a Binomial process, and the probability of a new arrival during each frame is p_A ;
- the probability of a service completion (and thus, a departure) during each frame is p_S , provided that there is at least one job in the system at the beginning of the frame;
- service times and interarrival times are *independent*;
- jobs are being serviced in the order of their arrival.

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Markov property

Obviously, a B1SQS is a Markov chain.

Since the probabilities p_A and p_S never change, it is also a homogeneous Markov chain.

The number of jobs in the system increases by 1 with every arrival and decreases by 1 with each departure.

Conditions of a Binomial process guarantee that at most one arrival and at most one departure may occur during each frame.

The states are $\{0, 1, \ldots\}$ (number of jobs in the system). Let us find the transition probabilities.

$$p_{00} = P(0 \text{ arrivals}) = 1 - p_A$$

 $p_{01} = P(1 \text{ arrival}) = p_A.$

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In general, for $i \ge 1$,

$$p_{i,i-1} = P(0 \text{ arrivals and } 1 \text{ departure}) = P(\{0 \text{ arrivals}\} \cap \{1 \text{ departure}\})$$
 $= (1 - p_A)p_S$
 $p_{i,i} = P((\{0 \text{ arrivals}\} \cap \{0 \text{ departures}\}) \cup (\{1 \text{ arrival}\} \cap \{1 \text{ departure}\}))$
 $= P(\{0 \text{ arrivals}\} \cap \{0 \text{ departures}\}) + P(\{1 \text{ arrival}\} \cap \{1 \text{ departure}\})$
 $= (1 - p_A)(1 - p_S) + p_A p_S$
 $p_{i,i+1} = P(\{1 \text{ arrival}\} \cap \{0 \text{ departures}\})$
 $= p_A(1 - p_S).$

All the other transition probabilities are 0, since the number of jobs cannot change by more than 1 in any single frame.



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So, the transition probability matrix is

$$P = \begin{bmatrix} 1 - p_A & p_A & 0 & \dots & 0 & \dots \\ (1 - p_A)p_S & (1 - p_A)(1 - p_S) & p_A(1 - p_S) & \dots & 0 & \dots \\ + p_Ap_S & p_A(1 - p_S) & \dots & 0 & \dots \\ 0 & (1 - p_A)p_S & (1 - p_A)(1 - p_S) & \dots & 0 & \dots \\ + p_Ap_S & \dots & 0 & \dots \\ \vdots & \vdots & \vdots & \ddots & \vdots \end{bmatrix},$$

an $\infty \times \infty$ tridiagonal matrix.



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Below, see the transition diagram.



Figure 1: Transition diagram for a B1SQS

The transition probability matrix may be used, for example, to simulate this queuing system and study its performance, as we did with general Markov chains.

One can also compute k-step transition probabilities and predict the load of a server or the length of a queue at any time in the future.

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Example 3.2.

Jobs (documents) are sent to a printer at the rate of 20 per hour. It takes an average of 40 seconds to print a document. Currently, the printer is printing a job, and there is another job stored in a queue. Assume a B1SQS with 20-second frames is modeling this printer.

- a) Compute the probability that the printer will be idle in 2 minutes.
- **b**) Find the expected total number of jobs in the system in 2 minutes.
- c) What is the expected length of the queue in 2 minutes?
- **d)** What is the expected waiting time for a document in 2 minutes?
- **e**) On average, how long does it take to get the printout of a document in 2 minutes?

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Solution.

First off, let us note that any printer represents a single-server queuing system, because it can process only one job at a time while other jobs are waiting in a queue.

Now, parameters are given in hours, in minutes and in seconds, so let us choose the "middle" one, i.e., express everything in minutes. We are given:

$$\lambda_A = 20 / \text{hour} = 1/3 / \text{minute, so}$$
 $\mu_A = 3 \text{ minutes,}$
 $\mu_S = 40 \text{ seconds} = 2/3 \text{ minutes, so}$
 $\lambda_S = 1/\mu_S = 3/2 / \text{minute,}$
 $\Delta = 20 \text{ seconds} = 1/3 \text{ minutes.}$

Then

$$p_A = \lambda_A \Delta = 1/9, 1 - p_A = 8/9,$$

 $p_S = \lambda_S \Delta = 1/2, 1 - p_S = 1/2.$

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The transition probabilities are

$$p_{00} = 1 - p_A = 8/9,$$

 $p_{01} = p_A = 1/9,$
 $p_{i,i-1} = (1 - p_A)p_S = 8/9 \cdot 1/2 = 4/9,$

$$p_{i,i} = (1 - p_A)(1 - p_S) + p_A p_S = 8/9 \cdot 1/2 + 1/9 \cdot 1/2 = 1/2,$$

 $p_{i,i+1} = p_A(1 - p_S) = 1/9 \cdot 1/2 = 1/18.$

Hence,

$$P = \begin{bmatrix} 8/9 & 1/9 & 0 & 0 & \dots \\ 4/9 & 1/2 & 1/18 & 0 & \dots \\ 0 & 4/9 & 1/2 & 1/18 & \dots \\ 0 & 0 & 4/9 & 1/2 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}.$$

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Now, in t=2 minutes, there are $n=\frac{t}{\Delta}=6$ frames. The distribution of X after 6 frames is

$$P_6 = P_0 \cdot P^6.$$

The initial distribution (2 jobs in the system) is

$$P_0 = [0 \ 0 \ 1 \ 0 \ \ldots].$$

Here is an interesting problem. How do we deal with matrix *P* that has infinitely many rows and columns?

Fortunately, we only need a small portion of this matrix.

In the course of 6 frames, the number of jobs in the system, X(t), can change by 6 at most (see Figure 1), i.e. it can reach a maximum of 8.

Thus, it is sufficient to consider the first 9 rows and 9 columns of P only, corresponding to states $\{0, 1, \dots, 8\}$.

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So, we consider P_0 (and P_6) as having length 9,

$$P_0 = [0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0]$$

and P a 9×9 matrix,

$$P \ = \ \begin{bmatrix} 8/9 & 1/9 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 4/9 & 1/2 & 1/18 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 4/9 & 1/2 & 1/18 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 4/9 & 1/2 & 1/18 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 4/9 & 1/2 & 1/18 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 4/9 & 1/2 & 1/18 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 4/9 & 1/2 & 1/18 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 4/9 & 1/2 & 1/18 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 4/9 & 1/2 & 1/18 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 4/9 & 1/2 & 1/18 \\ \end{bmatrix}$$

In 2 minutes (6 frames), the distribution will be

$$P_6 = P_0 \cdot P^6 = [0.6436 \ 0.25 \ 0.0799 \ 0.0218 \ 0.0041 \ 0.0005 \ 0 \ 0 \ 0].$$

Now we can answer all the questions.

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a) Compute the probability that the printer will be idle in 2 minutes.

The probability that the printer is idle after 2 minutes is the probability of 0 jobs in the system at that time, i.e.

$$P_6(0) = 0.6436.$$

b) Find the expected total number of jobs in the system in 2 minutes.

In 2 minutes, the total number of jobs in the system, X(2), has pdf

$$X\begin{pmatrix} 0 & \dots & 8 \\ P_6 & \end{pmatrix}$$

$$= \begin{pmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 0.6436 & 0.25 & 0.0799 & 0.0218 & 0.0041 & 0.0005 & 0 & 0 & 0 \end{pmatrix},$$

SO

$$E(X) = \sum_{k=0}^{8} kP_6(k) = 0.4944 \text{ jobs.}$$

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c) What is the expected length of the queue in 2 minutes?

Out of the X jobs in the system above, X_w jobs are waiting in a queue and X_s are being serviced. The expected length of the queue is the expected value

$$E(X_w) = E(X - X_s) = E(X) - E(X_s).$$

We have found E(X), so let us turn our attention to X_s .

Since the server (printer) can process at most 1 job at a time, X_s is either 0 or 1, i.e. it has a Bernoulli distribution.

With what parameter p?

The parameter is the probability of "success", in this case, the probability that the system is working, so, *not* idle:

$$p = P(\text{printer is busy}) = 1 - P(\text{printer is idle})$$

= $1 - P_6(0) = 1 - 0.6436 = 0.3564$.

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So the pdf of X_s is

$$X_s \left(\begin{array}{cc} 0 & 1 \\ 1 - p & p \end{array} \right), \ p = 0.3564$$

and its expected value

$$E(X_s) = p = 0.3564.$$

Then the expected queue length is

$$E(X_w) = E(X) - E(X_s) = 0.4944 - 0.3564 = 0.138 \text{ jobs.}$$

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d) What is the expected waiting time for a document in 2 minutes?

The expected waiting time for a document is E(W). By Little's Law, we have

$$E(W) = \frac{1}{\lambda_A} E(X_w) = \mu_A E(X_w)$$
= 3 · 0.138 minutes = 0.414 minutes
= 24.84 seconds.

e) On average, how long does it take to get the printout of a document in 2 minutes?

This is the expected total time the job spends in the system, i.e., the expected response time of a job, in 2 minutes. Again, by Little's Law, that number is

$$E(R) = \frac{1}{\lambda_A} E(X) = 3 \cdot 0.4944 \text{ minutes} = 1.4832 \text{ minutes}.$$



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Remark 3.3.

- **1.** A B1SQS is an irregular Markov chain. Any k-step transition probability matrix contains zeros because a k-step transition from 0 to k+1 is impossible. It requires at least k+1 arrivals, and this cannot happen by the conditions of the Binomial process of arrivals.
- **2.** However, without the Binomial counting process restrictions, it can be shown that any system whose service rate exceeds the arrival rate (i.e., jobs can be served faster than they arrive, so there is no overload),

$$\lambda_S > \lambda_A$$
,

does have a steady-state distribution. Its computation is possible, despite the infinite dimension of P, but a little complicated. Instead, we will compute the steady-state distribution of a continuous queuing process, obtained by letting the frame size $\Delta \to 0$.

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Systems with limited capacity

- as we have seen, the number of jobs in a B1SQS may potentially reach any value;
- however, in practice, many systems have limited resources for storing jobs;
- ➤ then, there is a maximum number of jobs *C* that can possibly be in the system simultaneously;
- this number is called capacity;
- as examples, consider people going to a restaurant, cars entering a parking lot, customers going into a bank, etc.

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How does the situation change for a queuing system with a limited capacity $C < \infty$?

Not much, but it does make a difference. Up until the capacity *C* is reached, the system operates as before.

Things change when X = C. At this time, the system is full, so it can accept new jobs into its queue only if some job departs. We have

$$p_{C,C-1} = P(0 \text{ arrivals } \cap 1 \text{ departure}) = (1 - p_A)p_S \text{ (as before)},$$
 $p_{C,C} = P((0 \text{ arrivals } \cap 0 \text{ departures}) \cup (1 \text{ arrival } \cap 1 \text{ departure})$
 $\cup (1 \text{ arrival } \cap 0 \text{ departures}))$
 $= (1 - p_A)(1 - p_S) + p_A p_S + p_A(1 - p_S)$
 $= 1 - (1 - p_A)p_S.$

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This Markov chain has states 0, 1, ..., C, its transition probability matrix is finite, and it is regular (any state can be reached in C steps).

The transition diagram for a system with limited capacity is given in Figure 2.

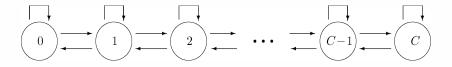


Figure 2: Transition diagram for a B1SQS with limited capacity C

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Example 3.4.

A customer service representative has a telephone with 2 lines, so she can talk to a customer while having another one "on hold". Suppose the representative gets an average of 10 calls per hour and the average phone conversation lasts 4 minutes. Assuming a B1SQS with 1-minute frames find the steady-state distribution and interpret it.

Solution.

Obviously, this is a system with limited capacity C = 2. When the capacity is reached and someone tries to call, (s)he will get a busy signal or voice mail. This Markov chain X(t) has 3 states, 0, 1, 2 and we have:

$$\lambda_A = 10 / \text{hour} = 1/6 / \text{minute},$$
 $\mu_S = 4 \text{ minutes}, \text{ so}$
 $\lambda_S = 1/\mu_S = 1/4 / \text{minute},$
 $\Delta = 1 \text{ minute}.$

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So,

$$p_A = \lambda_A \Delta = 1/6, 1 - p_A = 5/6,$$

 $p_S = \lambda_S \Delta = 1/4, 1 - p_S = 3/4.$

The transition probability matrix, of dimensions 3×3 , is

$$P = \begin{bmatrix} 1 - p_A & p_A & 0\\ (1 - p_A)p_S & (1 - p_A)(1 - p_S) + p_A p_S & p_A (1 - p_S)\\ 0 & (1 - p_A)p_S & 1 - (1 - p_A)p_S \end{bmatrix}$$
$$= \begin{bmatrix} 5/6 & 1/6 & 0\\ 5/24 & 2/3 & 1/8\\ 0 & 5/24 & 19/24 \end{bmatrix}.$$

The steady-state distribution is found, as usually, from the system

$$\begin{cases} \pi P &= \pi \\ \sum_{k=0}^{2} \pi_k &= 1. \end{cases}$$

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This leads to

$$\begin{cases} \pi_0 - \frac{5}{4}\pi_1 & = 0 \\ \frac{3}{5}\pi_1 - \pi_2 & = 0 \\ \pi_0 + \pi_1 + \pi_2 & = 1, \end{cases}$$

with solution

$$\pi_0 = \frac{25}{57} \approx 0.439,$$
 $\pi_1 = \frac{20}{57} \approx 0.351,$
 $\pi_2 = \frac{12}{57} \approx 0.21.$

Interpretation: 43.9% of the time the representative is not talking on the phone (and, implicitly, there is no one on hold), 35.1% of the time she talks to a customer, but the second line is open, and 21% of the time both lines are busy (one talking, one holding) and no new calls can get through.

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