

# Entropy Rates of a Stochastic Process

Best Achievable Data Compression

Radu Trîmbițaș

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## 1 Entropy Rates of a Stochastic Process

### Entropy rates

- The AEP states that  $nH(X)$  bits suffice on the average for  $n$  i.i.d. RVs
- What for dependent RVs?
- For stationary processes  $H(X_1, X_2, \dots, X_n)$  grows (asymptotically) linearly with  $n$  at a rate  $H(\mathcal{X})$  – the *entropy rate* of the process
- A *stochastic process*  $\{X_i\}_{i \in I}$  is an indexed sequence of random variables,  $X_i : S \rightarrow \mathcal{X}$  is a RV  $\forall i \in I$
- If  $I \subseteq \mathbb{N}$ ,  $\{X_1, X_2, \dots\}$  is a *discrete stochastic process*, called also a *discrete information source*.
- A discrete stochastic process is characterized by the joint probability mass function

$$P((X_1, X_2, \dots, X_n) = (x_1, x_2, \dots, x_n)) = p(x_1, x_2, \dots, x_n)$$

where  $(x_1, x_2, \dots, x_n) \in \mathcal{X}^n$ .

### 1.1 Markov chains

#### Markov chains

**Definition 1.** A stochastic process is said to be *stationary* if the joint distribution of any subset of the sequence of random variables is invariant with respect to shifts in the time index

$$P(X_1 = x_1, \dots, X_n = x_n) = P(X_{1+\ell} = x_1, \dots, X_{n+\ell} = x_n) \quad (1)$$

$\forall n, \ell$  and  $\forall x_1, x_2, \dots, x_n \in \mathcal{X}$ .

**Definition 2.** A discrete stochastic process  $\{X_1, X_2, \dots\}$  is said to be a *Markov chain* or *Markov process* if for  $n = 1, 2, \dots$

$$\begin{aligned} &P(X_{n+1} = x_{n+1} | X_n = x_n, X_{n-1} = x_{n-1}, \dots, X_1 = x_1) \\ &= P(X_{n+1} = x_{n+1} | X_n = x_n), \quad x_1, x_2, \dots, x_n, x_{n+1} \in \mathcal{X}. \end{aligned} \quad (2)$$

The joint pmf can be written as

$$p(x_1, x_2, \dots, x_n) = p(x_1)p(x_2|x_1) \dots p(x_n|x_{n-1}). \quad (3)$$

**Definition 3.** A Markov chain is said to be *time invariant* (*time homogeneous*) if the conditional probability  $p(x_{n+1}|x_n)$  does not depend on  $n$ ; that is for  $n = 1, 2, \dots$

$$P(X_{n+1} = b | X_n = a) = P(X_2 = b | X_1 = a), \quad \forall a, b \in \mathcal{X}. \quad (4)$$

*This property is assumed unless otherwise stated.*

- $\{X_i\}$  Markov chain,  $X_n$  is called the *state* at time  $n$
- A time-invariant Markov chain is characterized by its initial state and a *probability transition matrix*  $P = [P_{ij}]$ ,  $i, j = 1, \dots, m$ , where  $P_{ij} = P(X_{n+1} = j | X_n = i)$ .
- The Markov chain  $\{X_n\}$  is *irreducible* if it is possible to go from any state to another with a probability  $> 0$
- The Markov chain  $\{X_n\}$  is *aperiodic* if  $\forall$  state  $a$ , the possible times to go from  $a$  to  $a$  have highest common factor = 1.
- Markov chains are often described by a directed graph where the edges are labeled by the probability of going from one state to another.
- $p(x_n)$  - pmf of the random variable at time  $n$

$$p(x_{n+1}) = \sum_{x_n} p(x_n)P_{x_n x_{n+1}} \quad (5)$$

- A distribution on the states such that the distribution at time  $n + 1$  is the same as the distribution at time  $n$  is called a *stationary distribution* - so called because if the initial state of a Markov chain is drawn according to a stationary distribution, the Markov chain form a stationary process.
- If the finite-state Markov chain is irreducible and periodic, the stationary distribution is unique, and from any starting distribution, the distribution of  $X_n$  tends to a stationary distribution as  $n \rightarrow \infty$ .

*Example 4.* Consider a two-state Markov chain with a probability transition matrix

$$P = \begin{bmatrix} 1 - \alpha & \alpha \\ \beta & 1 - \beta \end{bmatrix}$$

(Figure 1)

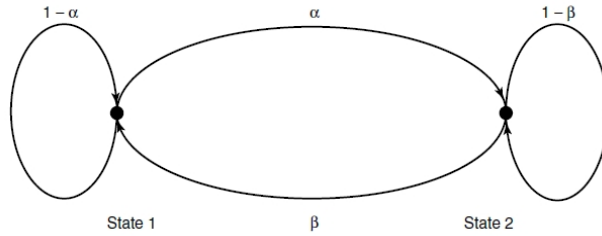


Figure 1: Two-state Markov chain

The stationary probability is the solution of  $\mu P = \mu$  or  $(I - P^T)\mu^T = 0$ . We add the condition  $\mu_1 + \mu_2 = 1$ .

The solution is

$$\mu_1 = \frac{\beta}{\alpha + \beta}, \quad \mu_2 = \frac{\alpha}{\alpha + \beta}.$$

Click here for a Maple solution [Markovex1.html](#). The entropy of  $X_n$  is

$$H(X_n) = H\left(\frac{\beta}{\alpha + \beta}, \frac{\alpha}{\alpha + \beta}\right).$$

## 1.2 Entropy rate

### Entropy rate

**Definition 5.** The *entropy rate* of a stochastic process  $\{X_i\}$  is defined by

$$H(\mathcal{X}) = \lim_{n \rightarrow \infty} \frac{1}{n} H(X_1, \dots, X_n) \quad (6)$$

when the limit exists.

#### Examples

1. Typewriter -  $m$  equally likely output letters; he/she can produce  $m^n$  sequences of length  $n$ , all of them equally likely.  $H(X_1, \dots, X_n) = \log m^n$ , and the entropy rate is  $H(\mathcal{X}) = \log m$  bits per symbol.
2.  $X_1, X_2, \dots$  i.i.d. RVs

$$H(\mathcal{X}) = \lim_{n \rightarrow \infty} \frac{H(X_1, \dots, X_n)}{n} = \lim_{n \rightarrow \infty} \frac{nH(X_1)}{n} = H(X_1).$$

3.  $X_1, X_2, \dots$  independent, but not identically distributed RVs

$$H(X_1, \dots, X_n) = \sum_{i=1}^n H(X_i)$$

It is possible that  $\frac{1}{n} \sum H(x_i)$  does not exist

**Definition 6.**

$$H'(\mathcal{X}) = \lim_{n \rightarrow \infty} H(X_n | X_{n-1}, X_{n-2}, \dots, X_1). \quad (7)$$

$H(\mathcal{X})$  is entropy per symbol of the  $n$  RVs;  $H'(\mathcal{X})$  is the conditional entropy of the last RV given the past.

For stationary processes both limits exist and are equal.

**Lemma 7.** For a stationary stochastic process,  $H(X_n | X_{n-1}, \dots, X_1)$  is nonincreasing in  $n$  and has a limit  $H'(\mathcal{X})$ .

*Proof.*

$$\begin{aligned} H(X_{n+1} | X_1, X_2, \dots, X_n) &\leq H(X_{n+1} | X_n, \dots, X_2) && \text{conditioning} \\ &= H(X_n | X_{n-1}, \dots, X_1). && \text{stationarity} \end{aligned}$$

$(H(X_n | X_{n-1}, \dots, X_1))_n$  is decreasing and nonnegative, so it has a limit  $H'(\mathcal{X})$ .  $\square$

**Lemma 8 (Cesáro).** If  $a_n \rightarrow a$  and  $b_n = \frac{1}{n} \sum_{i=1}^n a_i$  then  $b_n \rightarrow a$ .

**Theorem 9.** For a stationary stochastic process  $H(\mathcal{X})$  (given by (6)) and  $H'(\mathcal{X})$  (given by (7)) exist and

$$H(\mathcal{X}) = H'(\mathcal{X}). \quad (8)$$

*Proof.* By the chain rule,

$$\frac{H(X_1, \dots, X_n)}{n} = \frac{1}{n} \sum_{i=1}^n H(X_i | X_{i-1}, \dots, X_1).$$

But,

$$\begin{aligned} H(\mathcal{X}) &= \lim_{n \rightarrow \infty} \frac{H(X_1, \dots, X_n)}{n} \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n H(X_i | X_{i-1}, \dots, X_1) \\ &= \lim_{n \rightarrow \infty} H(X_n | X_{n-1}, \dots, X_1) && \text{(Lemma 8)} \\ &= H'(\mathcal{X}) && \text{(Lemma 7)} \end{aligned}$$

$\square$

### 1.3 Entropy rate for Markov chain

#### Entropy rate for Markov chain

- For a stationary Markov chain, the entropy rate is given by

$$\begin{aligned} H(\mathcal{X}) = H'(\mathcal{X}) &= \lim H(X_n | X_{n-1}, \dots, X_1) = \lim H(X_n | X_{n-1}) \\ &= H(X_2 | X_1), \end{aligned} \quad (9)$$

where the conditional entropy is calculated using the given stationary distribution.

- The stationary distribution  $\mu$  is the solution of the equations

$$\mu_j = \sum_i \mu_i P_{ij}, \quad \forall j.$$

- Expression of conditional entropy:

**Theorem 10.**  $\{X_i\}$  stationary Markov chain with stationary distribution  $\mu$  and transition matrix  $P$ . Let  $X_1 \sim \mu$ . then the entropy rate is

$$H(\mathcal{X}) = - \sum_i \sum_j \mu_i P_{ij} \log P_{ij}. \quad (10)$$

*Proof.*  $H(\mathcal{X}) = H(X_2|X_1) = \sum_i \mu_i \left( - \sum_j P_{ij} \log P_{ij} \right)$ . □

*Example 11* (Two-state Markov chain). The entropy rate of the two-state Markov chain in Figure 1 is

$$H(\mathcal{X}) = H(X_2|X_1) = \frac{\beta}{\alpha + \beta} H(\alpha) + \frac{\alpha}{\alpha + \beta} H(\beta).$$

**Remark.** If the Markov chain is irreducible and aperiodic, it has a unique stationary distribution on the states, and any initial distribution tends to the stationary distribution as  $n \rightarrow \infty$ . In this case, even though the initial distribution is not the stationary distribution, the entropy rate, which is defined in terms of long-term behavior, is  $H(\mathcal{X})$ , as defined in (9) and (10).

## 1.4 Functions of Markov chains

### Functions of Markov chains

- $X_1, X_2, \dots, X_n, \dots$  stationary Markov chain,  $Y_i = \phi(X_i)$ ,  $H(\mathcal{Y}) = ?$
- in many cases  $Y_1, Y_2, \dots, Y_n, \dots$  is not a Markov chain, but it is stationary
- lower bound

**Lemma 12.**

$$H(Y_n | Y_{n-1}, \dots, Y_2, X_1) \leq H(\mathcal{Y}). \quad (11)$$

*Proof.* For  $k = 1, 2, \dots$

$$\begin{aligned} H(Y_n | Y_{n-1}, \dots, Y_2, X_1) &\stackrel{(a)}{=} H(Y_n | Y_{n-1}, \dots, Y_2, Y_1, X_1) \\ &\stackrel{(b)}{=} H(Y_n | Y_{n-1}, \dots, Y_2, Y_1, X_1, X_0, X_{-1}, \dots, X_{-k}) \\ &\stackrel{(c)}{=} H(Y_n | Y_{n-1}, \dots, Y_2, Y_1, X_1, X_0, X_{-1}, \dots, \\ &\quad X_{-k}, Y_0, \dots, Y_{-k}) \\ &\stackrel{(d)}{\leq} H(Y_n | Y_{n-1}, \dots, Y_2, Y_1, Y_0, \dots, Y_{-k}) \\ &\stackrel{(e)}{=} H(Y_{n+k+1} | Y_{n+k}, \dots, Y_1), \end{aligned}$$

(a) follows from the fact that  $Y_1 = \phi(X_1)$ , (b) from the Markovity, (c) from  $Y_i = \phi(X_i)$ , (d) conditioning reduces entropy, (e) stationarity.  $\square$

*Proof - continuation.* Since inequality is true for all  $k$ , in the limit

$$\begin{aligned} H(Y_n|Y_{n-1}, \dots, Y_2, X_1) &\leq \lim_k H(Y_{n+k+1}|Y_{n+k}, \dots, Y_1) \\ &= H(\mathcal{Y}). \end{aligned}$$

$\square$

**Lemma 13.**

$$H(Y_n|Y_{n-1}, \dots, Y_2, X_1) - H(Y_n|Y_{n-1}, \dots, Y_2, Y_1, X_1) \rightarrow 0. \quad (12)$$

*Proof.* Expression of interval length:

$$\begin{aligned} H(Y_n|Y_{n-1}, \dots, Y_2, X_1) - H(Y_n|Y_{n-1}, \dots, Y_2, Y_1, X_1) \\ = I(X_1; Y_n|Y_{n-1}, \dots, Y_1). \end{aligned}$$

By properties of mutual information,

$$I(X_1; Y_1, \dots, Y_n) \leq H(X_1),$$

and  $I(X_1; Y_1, \dots, Y_n)$  increases with  $n$ . Thus,  $\lim I(X_1; Y_1, \dots, Y_n)$  exists and

$$\lim_{n \rightarrow \infty} I(X_1; Y_1, \dots, Y_n) \leq H(X_1).$$

$\square$

*Proof - continuation.* By the chain rule

$$\begin{aligned} H(X_1) &\geq \lim_{n \rightarrow \infty} I(X_1; Y_1, \dots, Y_n) \\ &= \lim_{n \rightarrow \infty} \sum_{i=1}^n I(X_1; Y_i|Y_{i-1}, \dots, Y_1) \\ &= \sum_{i=1}^{\infty} I(X_1; Y_i|Y_{i-1}, \dots, Y_1) \end{aligned}$$

The general term of the series must tend to 0

$$\lim I(X_1; Y_n|Y_{n-1}, \dots, Y_1) = 0.$$

$\square$

The last two lemmas imply

**Theorem 14.**  $X_1, X_2, \dots, X_n, \dots$  stationary Markov chain,  $Y_i = \phi(X_i)$

$$H(Y_n|Y_{n-1}, \dots, Y_1, X_1) \leq H(\mathcal{Y}) \leq H(Y_n|Y_{n-1}, \dots, Y_1) \quad (13)$$

and

$$\lim H(Y_n|Y_{n-1}, \dots, Y_1, X_1) = H(\mathcal{Y}) = \lim H(Y_n|Y_{n-1}, \dots, Y_1) \quad (14)$$

## Hidden Markov models

- We could consider  $Y_i$  to be a stochastic function of  $X_i$
- $X_1, X_2, \dots, X_n, \dots$  stationary Markov chain,  $Y_1, Y_2, \dots, Y_n, \dots$  a new process where  $Y_i$  is drawn according to  $p(y_i|x_i)$ , conditionally independent of all the other  $X_j, j \neq i$

$$p(x^n, y^n) = p(x_1) \prod_{i=1}^{n-1} p(x_{i+1}|x_i) \prod_{i=1}^n p(y_i|x_i).$$

- $Y_1, Y_2, \dots, Y_n, \dots$  is called a hidden Markov model (HMM)
- Applied to speech recognition, handwriting recognition, and so on.
- The same argument as for functions of Markov chain works for HMMs.

## References

## References

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- [3] Robert M. Gray, *Entropy and Information Theory*, Springer, 2009