

3.2. Kernels, cokernels, products and coproducts

In the sequel \mathcal{C} will be a preadditive category.

A zero object in \mathcal{C} is an object $0 \in \mathcal{C}$ such that $\mathcal{C}(X, 0) = \{0\}$ and $\mathcal{C}(0, X) = \{0\}$ for every $X \in \mathcal{C}$. Any two zero objects are isomorphic, therefore they will be identified to a single zero object $0 \in \mathcal{C}$. If \mathcal{C} has a zero object 0 exists, then the neutral element of the group $\mathcal{C}(X, X)$ is identical to the composed morphism $X \xrightarrow{\text{id}} 0 \rightarrow X$, as follows from the bilinearity of the composition of maps.

Example 3.2.1 The module with a single element $0 = \{0\}$ is a zero object in the category $R\text{-Mod}$.

Exercise 1 Show that in an additive category \mathcal{C} , a monomorphism $\alpha: X' \rightarrow X$

is a monomorphism (resp. epimorphism) iff for every $\beta: X'' \rightarrow X'$ (resp. $\gamma: X \rightarrow X''$) we have $\alpha\beta = 0$ (resp. $\beta\alpha = 0$) implies $\beta = 0$.

Two monomorphisms $\alpha: X \rightarrow Y$ and $\alpha': X' \rightarrow Y$ are called equivalent if there is an isomorphism $\gamma: X \rightarrow X'$ such that $\alpha'\gamma = \alpha$.

Exercise 2 Show that the relation defined above is actually an equivalence relation on the class of all monomorphisms with the same codomain $Y \in \mathcal{C}$.

An equivalence class of monomorphisms with the codomain Y is called a subobject of Y . By abuse of notation we will identify a subobject to a monomorphism representing it. If $X \xrightarrow{\alpha} Y$ and $X' \xrightarrow{\alpha'} Y$ are two subobjects of Y we write $X \leq X'$ if there is $\gamma: X \rightarrow X'$ such that $\alpha'\gamma = \alpha$.

Exercise 3 If $X \xrightarrow{\alpha} Y$ and $X' \xrightarrow{\alpha'} Y$ are subobjects and $X \leq X'$ then show that the morphism $\gamma: X \rightarrow X'$ for which $\alpha'\gamma = \alpha$ is a monomorphism. Show also that the relation " \leq " is reflexive and transitive on the class of all monomorphisms with the codomain Y . Consequently it induces an order on the class of all subobjects of Y (that is a relation which is reflexive, transitive and antisymmetric).

The dual notion of the notion of "subobject" is the so called quotient object. That is a quotient object is an equivalence class of epimorphisms with the same domain, where $X \xrightarrow{\beta} Y$, $X \xrightarrow{\beta'} Y'$ are equiv. if there exists an isomorphism $\gamma: Y \rightarrow Y'$ such that $\gamma\beta = \beta'$.

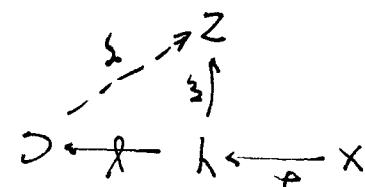
Lemma 3.2.1 If $\alpha: X \rightarrow Y$ is a morphism of \mathcal{A} . Then $\alpha = 0$ or there is a unique morphism $\beta: X \rightarrow Y$ such that $\alpha \circ \beta = 0$.

Proof. Let $\alpha: X \rightarrow Y$ be the kernel of $\beta: Y \rightarrow Z$, and let $\gamma: Y \rightarrow C$ be the cokernel of α .

Proposition 3.2.4 If $\alpha: X \rightarrow Y$ is the kernel of some morphism $\beta: X \rightarrow C$ and $\beta: Y \rightarrow Z$ is the cokernel of the same morphism $\gamma: C \rightarrow Z$, then $\alpha = \beta \circ \gamma$.

By definition it follows that the cokernel

$$\text{coker } \beta = Y / \text{ker } \beta = Y / \text{ker } (\text{coker } \alpha) = Y / \text{ker } \alpha = Y / \text{ker } (\text{coker } \beta).$$



If $\alpha = 0$ and for every $\beta: Y \rightarrow Z$ such that $\beta \circ \alpha = 0$ then by definition. This means $\beta: Y \rightarrow C$ is called the cokernel of $\alpha: X \rightarrow Y$.

The notion of cokernel is defined from the one of the kernel

(and exercise)

Proposition 3.2.3 A morphism is a monomorphism if $\text{ker } \alpha = 0$.

About the kernel of α if it exists. We denote the kernel by $\text{ker } \alpha$.

Remark That because the unit of the field is $1_{\mathcal{A}}$. 3.2.2. we may note

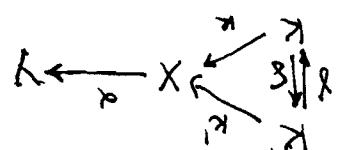
is a monomorphism and the above diagram shows that two kernels

of $\alpha: X \rightarrow Y$ if $\text{ker } \alpha$ by $\text{ker } \alpha = 0$, $\text{ker } \alpha = \text{ker } \beta$ to α

we observe by unitarity stated in the definition

$$g \circ f = k_1 = k_2 = k_3 = k_4$$

$$\text{thus } k_3 g = k_4 f = k_1 = k_2 \text{ and}$$



morphisms $g: K \rightarrow X$ and $f: K \rightarrow Y$, $k_1 = k_2$ and $k_3 = k_4$

as the kernels of α to form by definition of the kernel from our

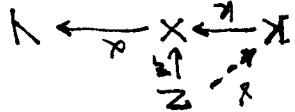
Proposition The unitarity of the kernel is standard: If $K \xrightarrow{\alpha} X$

such that $\text{ker } \alpha = K$ and $\text{coker } \alpha = X$.

Proposition 3.2.2 The kernel of a morphism is a unique

one since:

Remember that the definition of a kernel is an universal property. Therefore



$$\alpha = 0 \text{ where } \beta = 0 \text{ and for every } \beta: X \rightarrow Z \text{ such that } \beta \circ \alpha = 0$$

such that $\alpha \circ \beta = 0$ and $\beta: X \rightarrow Z$ such that $\beta \circ \alpha = 0$.

A kernel of a morphism $\alpha: X \rightarrow Y$ is the unique monomorphism $\alpha: K \rightarrow X$

We have that $\pi_1^* \alpha \beta = \pi_2^* \alpha \beta$ or $\alpha \beta = \pi_1^* \alpha \beta = \alpha \beta$ for all α, β . Thus π_1^* and π_2^* are unitary morphisms of the products $\mathbb{H}X$: π_1^* is surjective!

Example 3.2.2 If $\alpha: X \rightarrow Y$ is an isomorphism, then so is π_1^* . To the product map is induced as a functor $\mathcal{G} \rightarrow \mathcal{G}$.

When the weak cell morphisms are the canonical morphisms of the product

$$\begin{array}{ccc} X_{\mathbb{H}} & \dashrightarrow & Y_{\mathbb{H}} \\ \downarrow & \alpha & \downarrow \\ X & \xleftarrow{\alpha} & Y \end{array}$$

we obtain a mapping commutative the diagram

through α , we define the product $\alpha^* \mathbb{H} \alpha$ as being the unique family $\alpha^*: X_i \rightarrow Y_i$ ($i \in I$) to morphisms in a category with a couniversal property

We denote by $\mathbb{H}X$ the product of the form $\mathbb{H}(X)$ (X is a collection

Exercise 3 Show that the product is unique up to a natural isomorphism.

(i.e. The diagram is commutative)

It is a universal morphism with $\mathbb{H}f = f^* \mathbb{H}f$ such that $f^* \mathbb{H}f = f$ and $\mathbb{H}f \circ f^* \mathbb{H}f = f$ ($f \in I$). This means that for each pair $Y \in \mathcal{G}$ and morphisms $\mathbb{H}f: Y \rightarrow X$ ($f \in I$) there is a unique $X \in \mathcal{G}$ together with morphisms $\mathbb{H}f: Y \rightarrow X$ ($f \in I$) such that (X, f) is a pair satisfying $\mathbb{H}f \circ f^* \mathbb{H}f = f$. A product of this form

The colimit of α is the profibration $Y \rightarrow Y/\text{im } \alpha$.

$\text{colim } \alpha = \{x \in X / \alpha(x) = 0\}$ (together with the inclusion $\text{colim } \alpha \hookrightarrow X$)

where $\alpha: X \rightarrow Y$ is the forgetful bimodule $\alpha = \text{colim } \alpha / \alpha(0) = 0$

Example 3.2.5 In the category $R\text{-Mod}$ the limit of our functors

that $\alpha = \text{colim } \alpha$.

is a square $b: \mathbb{H}X \rightarrow X$ s.t. $\alpha \circ b = f$. This means

$b \circ f = g = g \circ b = 0$ and thus $\alpha = \text{colim } \alpha$ is the object

for $f: U \rightarrow Y$ and $b: U \rightarrow X$ it follows

$$\begin{array}{ccccc} & & 2 & \rightarrow & \\ & & \downarrow & & \\ & & 3 & \rightarrow & \\ & & \downarrow & & \\ C & \leftarrow & X & \leftarrow & Y \\ & & \uparrow & & \\ & & U & \leftarrow & V \end{array}$$

Dualizing the notion of product we obtain the coproduct which is also called direct sum. Therefore an object $X \in \mathcal{C}$ together with a family of morphisms $f_i : X_i \rightarrow X$ ($i \in I$) is called the direct sum of the family of objects $X_i \in \mathcal{C}$ ($i \in I$) if for every family of morphisms $\xi_i : X_i \rightarrow Y$ there is a unique morphism $\xi : X \rightarrow Y$ s.t. $\xi_i = \xi \circ f_i$, ($i \in I$).

$$\begin{array}{ccc} X_i & \xrightarrow{f_i} & * \\ & \searrow \xi & \\ & \xi_i & \rightarrow Y \end{array}$$

We denote $\bigoplus_{i \in I} X_i$ the direct sum and we call f_i canonical injections of the direct sum.

Exercise 6 Show that the canonical projections of the product are epimorphisms and the canonical injections of the direct sum are monomorphisms.

Exercise 7 Show that if $\{X_i \in \mathcal{C} \mid i \in I\}$ ~~with the \oplus~~ has the product for all $X \in \mathcal{C}$, in \mathcal{C} then it holds $\mathcal{C}(X, \prod_i X_i) \cong \prod_i \mathcal{C}(X, X_i)$ where the second product is taken in the category \mathbf{Ab} of all abelian groups. Conversely if $X \in \mathcal{C}$ is an object satisfying $\mathcal{C}(Y, X) \cong \prod_{i \in I} \mathcal{C}(Y, X_i)$ for all $Y \in \mathcal{C}$ then $X = \prod_{i \in I} X_i$. Dually the direct sum is characterized by the formula: $\mathcal{C}(\bigoplus X_i, X) \cong \prod_i \mathcal{C}(X_i, X)$ for all $X \in \mathcal{C}$.

Example 3.2.7 In category $R\text{-Mod}$, the products and exist products and direct sums for all families

Example 3.2.7a) In category $R\text{-Mod}$, every family of objects (R -modules) has product and direct sum, as we have seen in Ch. 1, Section 3.
 b) In the category Set , every family of objects has product and direct sum, namely the cartesian product resp. the disjoint union. (Exercise 8).

Note that, in general neither the product nor the direct sum of a family of objects in an arbitrary category ~~need not~~ do need not exist. Our presentation is oriented to module categories where they are defined. As a first step, ~~the~~ the reader unfamiliar with the language of categories may see all our constructions in categories as another approach of modules, with the above ~~warning~~ that not all is valid for arbitrary categories. Warning

Let $\{X_i\}_{i \in I}$ be a family of objects in \mathcal{C} such that both $\prod X_i$ and $\bigoplus X_i$ exist. For any $i, j \in I$ we define the morphisms $s_{ij} : C_j \rightarrow C_i$ as

$$s_{ij} = \begin{cases} 1 & \text{if } i=j \\ 0 & \text{if } i \neq j \end{cases} \quad (\text{the Kronecker delta})$$

From the definition of the product (direct sum), for every $j \in I$ there is a unique morphism $f'_j : X_j \rightarrow \prod X_i$ (resp. $\pi'_j : \bigoplus X_i \rightarrow X_j$) such that $\pi'_i f'_j = s_{ij}$ (resp. $\pi'_j f_i = s_{ji}$) for all $i \in I$. (Here we denoted π'_i and f'_i (resp.) the canonical projection resp. inj. of the product resp. direct sum.) In the case when I is finite, say $I = \{1, 2, \dots, n\}$, we get

Lemma 3.2.8 a) $s'_1 \pi'_1 + \dots + s'_n \pi'_n = 1_{\prod X_i}$

$$\text{b) } s_1 \pi'_1 + \dots + s_n \pi'_n = 1_{\bigoplus X_i}$$

(the notations are those used above).

Proof a) For all $i \in I$ we have $\pi'_i (s'_1 \pi'_1 + \dots + s'_n \pi'_n) = s'_{ii} \pi'_i + \dots + s'_{in} \pi'_n = \pi'_i = \pi'_i \cdot 1$ so $s'_1 \pi'_1 + \dots + s'_n \pi'_n = 1$ by the unicity of the factorization through the product.

3) Exercise 9

Conversely we have:

Proposition 3.2.9 Suppose there are given morphisms $f_i : X_i \rightarrow X$ and $\pi_i : X \rightarrow X_i$, with $1 \leq i \leq n$, such that

$$\pi_i f_j = s_{ij} \quad (\forall i, j \leq n) \text{ and } f_1 \pi_1 + \dots + f_n \pi_n = 1_X$$

Then X is the product and the coproduct of objects $(X_i)_{i \in \overline{1, n}}$ with the projections π_i and the injections f_i .

Proof First we shall show that X is the coproduct. Let $\xi_i : X_i \rightarrow Y$, $1 \leq i \leq n$. Define $\xi = \xi_1 \pi_1 + \dots + \xi_n \pi_n : X \rightarrow Y$. Then $\xi \xi_i = (\xi_1 \pi_1 + \dots + \xi_n \pi_n) f_i = \xi_i f_i + \dots + \xi_n f_n = \xi_i$ ($1 \leq i \leq n$). Moreover if $\eta : X \rightarrow Y$ such that $\eta f_i = \xi_i$ ($1 \leq i \leq n$) then $\eta = \xi \cdot 1_X = \xi (\xi_1 \pi_1 + \dots + \xi_n \pi_n) = \xi_1 \pi_1 + \dots + \xi_n \pi_n = \xi_1 \pi_1 + \dots + \xi_n \pi_n = \xi$ so ξ is unique. Thus $X = \bigoplus_{i=1}^n X_i$. To show that $X = \prod_{i=1}^n X_i$ is Exercise 10.

The above result tells us that ~~we need to distinguish~~ between products and direct sums over finite index sets coincide. We use notation $X_1 \times X_2 \times \dots \times X_n$ or $- \otimes - X_1 \oplus X_2 \oplus \dots \oplus X_n$.

Exercise 13 Let $\alpha: X \rightarrow Y$ be a morphism in a preadditive category. Show that $X = Y \oplus \ker \alpha$.
 It's clear that $\alpha \circ h = 0$ for all $h \in \text{Hom}(X, \ker \alpha)$. If $h \in \text{Hom}(Y, \ker \alpha)$, then $h \circ \alpha = 0$.
 This shows that $\ker \alpha$ is a subobject of Y in a preadditive category.

Exercise 12 Suppose that $\text{Hom}(X, Y) = \{x \in G \mid T(x) = \{x \in G \mid f(x) = g\}\}$. Then f and g in $A \oplus B$.
 That is, f and g are morphisms from X to $A \oplus B$ such that $f \circ h = g \circ h$ for all $h \in \text{Hom}(X, A \oplus B)$.
 This means that $f(h) = g(h)$. Recall that $f(h) = \{x \in G \mid f(x) = h\}$.
 Now $f(h) = g(h)$ means that $x \in f(h) \iff x \in g(h)$.
 Since $x \in f(h) \iff f(x) = h \iff f(x) = g(x)$, it follows that $x \in g(h)$.
 Thus $f(h) = g(h)$ for all $h \in \text{Hom}(X, A \oplus B)$.
 And if \oplus is also a direct sum in G , then \oplus is a direct sum in G .
 This shows that \oplus is a direct sum in G .

Exercise 11. Consider the full subcategory \mathcal{D} of the category \mathcal{C} to which \mathcal{D} is a full subcategory of \mathcal{C} .
 Show that $\mathcal{D} \oplus \mathcal{D} \cong \mathcal{D}$.

To show that $\mathcal{D} \oplus \mathcal{D} \cong \mathcal{D}$, we need to find a morphism $T: \mathcal{D} \oplus \mathcal{D} \rightarrow \mathcal{D}$ such that $T \circ (f \oplus g) = f$ and $T \circ (g \oplus f) = g$.
 Let $f, g \in \mathcal{D}$. We want to find $T(f, g)$ such that $T(f, g) \in \mathcal{D}$ and $T(f, g) \oplus T(f, g) = (f, g)$.
 This means that $T(f, g) \in \mathcal{D}$ and $T(f, g) \oplus T(f, g) = (f, g)$.
 Let $T(f, g) = (f \oplus g, f \oplus g)$. Then $T(f, g) \in \mathcal{D}$ because $f, g \in \mathcal{D}$.
 Now $T(f, g) \oplus T(f, g) = (f \oplus g, f \oplus g) \oplus (f \oplus g, f \oplus g) = (f \oplus g \oplus f \oplus g, f \oplus g \oplus f \oplus g) = (f, g)$.
 Therefore $T(f, g) \oplus T(f, g) = (f, g)$.
 Hence $T: \mathcal{D} \oplus \mathcal{D} \rightarrow \mathcal{D}$ is a morphism such that $T \circ (f \oplus g) = f$ and $T \circ (g \oplus f) = g$.
 Therefore $\mathcal{D} \oplus \mathcal{D} \cong \mathcal{D}$.

Proof that T is additive. Let $f, g \in \mathcal{D}$ and $\alpha, \beta \in \text{Hom}(X, \mathcal{D})$.
 We want to show that $T(\alpha + \beta, f + g) = T(\alpha, f) + T(\beta, g)$.
 By definition $T(\alpha + \beta, f + g) = (f + g) \oplus (\alpha + \beta)$.
 Now $(f + g) \oplus (\alpha + \beta) = f \oplus g \oplus \alpha \oplus \beta$.
 Since \oplus is a direct sum in \mathcal{D} , we have $f \oplus g = f$ and $\alpha \oplus \beta = \alpha$.
 Therefore $T(\alpha + \beta, f + g) = f \oplus g \oplus \alpha \oplus \beta = f \oplus \alpha \oplus g \oplus \beta = T(\alpha, f) + T(\beta, g)$.
 Hence T is additive.