

To my father, as a token  
of gratitude

## A VARIETY OF ASSOCIATIVE STRUCTURE WITH ONE-SIDED ZERO ELEMENTS IN AUTONOMOUS CATEGORIES

BY

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Using the "lifting" of a well-known bijection in *Ens* and the definition of the multiplicative systems in an autonomous category, a variety in the Frölich sense is found. The variety contains multiplicative systems "lifted" over associative structures of one-sided zero elements.

1. *Introduction.* It is a well known fact that in the category *Ens*, the following bijection holds

$$\alpha_{AB} : \text{Hom}(A, B) \cong \prod_{\alpha \in A} B$$

where  $A, B$  are arbitrary sets and  $\alpha_{AB}(f) = \{f(\alpha)\}_{\alpha \in A}$ .

In categories with richer structure, like  $Ab$  or  $Mod_R$ ,  $\alpha$  remains only an injective homomorphism, i.e. a monomorphism in these categories.

Let  $\mathcal{a}$  be a closed category in the sense of (2). We shall denote by  $\langle A, B \rangle$  the values of the lifted hom-functor  $\text{hom}_{\mathcal{a}} : \mathcal{a}^{\text{op}} \times \mathcal{a} \rightarrow \mathcal{a}$ , and by  $|A|$  the value of the set-valued functor of subjacency  $|-| : \mathcal{a} \rightarrow \text{Ens}$ , on objects  $A, B$  in  $\mathcal{a}$ .

In (4), if  $A, B$  are objects in  $\mathcal{a}$  and  $\alpha \in |A|$ , the "evaluation in  $\alpha$ "-morphism is constructed as follows

$$ev_{\alpha} : \langle A, B \rangle \rightarrow B$$

is the image of  $\alpha$  by the subjacent map of the  $\mathcal{a}$ -morphism

$$A \rightarrow \langle \langle A, B \rangle, B \rangle$$

corresponding to the identity of  $\langle A, B \rangle$  in the bijection

$$|\langle \langle A, B \rangle, \langle A, B \rangle \rangle| \cong |A, \langle \langle A, B \rangle, B \rangle|.$$

One can easily show that for every  $f \in |(A, B)|$  we have

$$|ev_\alpha|(f) = |f|(\alpha).$$

If  $\mathcal{a}$  has products, the family  $\{ev_\alpha : (A, B) \rightarrow A\}_{\alpha \in |A|}$  induces by factorization a morphism  $\alpha_{AB} : (A, B) \rightarrow \prod_{\alpha \in |A|} B$ .

Assuming that the functor of subjacency is faithful and preserves products (assumptions that hold in most of the concrete categories) we can trivially show that  $\alpha$  is monic and we have

$$|\alpha|(f) = \{|f|(\alpha)\}_{\alpha \in |A|}.$$

The commutativity of the two following diagrams shows that  $\alpha_{AB}$  is natural in his both arguments :

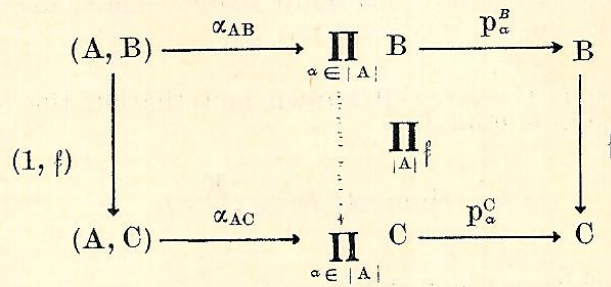


Diagram 1

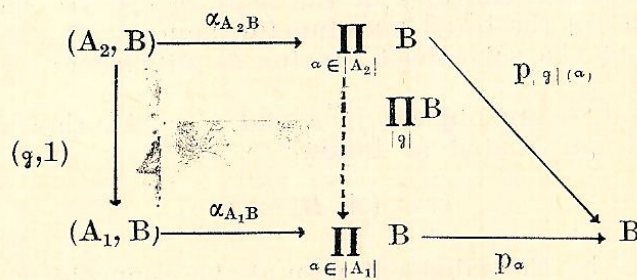


Diagram 2

The commutativity holds because equalities like

$$ev_\alpha^C \cdot (1, f) = f \cdot ev_\alpha^B, \quad ev_\alpha \cdot (g, 1) = ev_{|g|(\alpha)}$$

for any  $f : B \rightarrow C$  and  $g : A_1 \rightarrow A_2$  are easily established at the subjacent level (this is sufficient because of the faithfulness of  $|-|$ ).



2. DEFINITIONS. Let  $\mathfrak{a}$  be a closed category. We define, as in (1) the category  $\mathcal{M}$  of *multiplicative systems* in  $\mathfrak{a}$ ; that is, the objects of  $\mathcal{M}$  are pairs  $(x, \pi)$  where  $x \in \mathfrak{a}$  and  $\pi : x \rightarrow (x, x)$  is an arrow of  $\mathfrak{a}$ , and the morphisms are the arrows  $f : (x, \pi) \rightarrow (x', \pi')$  which are just arrows  $f : x \rightarrow x'$  in  $\mathfrak{a}$  such that the following pentagon commutes.

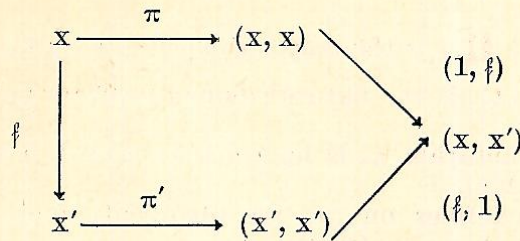


Diagram 3

It is easy to see that every multiplicative system defines an operation on  $|x|$  in a trivial way : for  $\alpha, \beta \in |x|$ ,  $\alpha \circ \beta = ||\pi|(\alpha)|(\beta)$ .

First, we shall show that if  $\mathfrak{a}$  is a symmetric monoidal closed category with products, then  $\mathcal{M}$  has products.

Let  $\{(x_i, \pi_i)\}_{i \in I}$  be a family of multiplicative systems. Denoting by  $P$  the product in  $\mathfrak{a}$   $\prod_i x_i$ , we are looking for a morphism,  $\pi : P \rightarrow (P, P)$  induced by the family  $\{\pi_i\}$ .

We consider the family

$$\left\{ P \xrightarrow{pr_i} x_i \xrightarrow{\pi_i} (x_i, x_i) \xrightarrow{(pr_i, 1)} (P, x_i) \right\}_{i \in I}.$$

$\mathfrak{a}$  being symmetric monoidal closed, the functor  $\mathfrak{a}(P, -) : \mathfrak{a} \rightarrow \mathfrak{a}$  has a left (strong) adjoint namely  $- \otimes P : \mathfrak{a} \rightarrow \mathfrak{a}$ , and therefore preserve limits and in particular, products.

It follows that  $\prod_{i \in I} (P, x_i) \cong (P, \prod_i x_i) = (P, P)$ , and so the above family gives by factorization the morphism  $\pi : P \rightarrow (P, P)$  we have been looking for. In this way,  $pr_i : (P, \pi) \rightarrow (x_i, \pi_i)$  is a morphism in  $\mathcal{M}$ .

*Remark.* The multiplicative system  $(P, \pi)$  is canonic enough; in fact, one can show that the "operation"  $\pi$  is induced by  $\{\pi_i\}$  "componentwise".

Recall from (3) that a *variety*  $\mathcal{B}$  in a category  $\mathcal{C}$  is a full subcategory of  $\mathcal{C}$ , satisfying the following axioms :

- (i) If  $f : A \rightarrow B$  is an epimorphism and  $A \in \mathcal{B}$  then  $B \in \mathcal{B}$
- (ii) If  $f : A \rightarrow B$  is a monomorphism and  $B \in \mathcal{B}$  then  $A \in \mathcal{B}$
- (iii) If  $(A_i)_{i \in I}$  is an indexed set of objects,  $(A_i) \subset \mathcal{B}$  then

$$\prod A_i \in \mathcal{B}.$$



3. *The variety.* Let  $\tilde{\mathcal{M}}$  be the full subcategory of  $\mathcal{M}$  whose objects are multiplicative systems  $(x, \pi_x)$  where  $\pi_x$  is such, that the two morphisms

$$\begin{array}{ccc} x & \xrightarrow{\Delta_x} & \prod_{|x|} x \\ \pi_x \searrow & & \nearrow \alpha_{xx} \\ & (x, x) & \end{array}$$

are equal.  $\Delta_x : x \rightarrow \prod_{|x|} x$  is the *diagonal morphism* induced by the family  $\{id_x : x \rightarrow x\}_{|x|}$  and  $\alpha_{xx}$  is the natural monomorphism defined in the first section.

*Remark.* The notation  $\pi_x$  is legitimate : if such a morphism exists, it is unique ( $\alpha$  monic).

*Remark.*  $|\pi_x|$  defines on  $|x|$  an operation such that for every  $i, j \in |x|$  we have  $i \circ j = |\pi_x|(i)(j) = i$ .

Thus, all the elements of  $|x|$  are all left-zero or all right-zero elements for the operation  $|\pi_x|$ . It is easy to see that this operation is associative, but not commutative.

**THEOREM.** *Let  $\mathcal{M}$  be the category of multiplicative systems on an autonomous category  $\alpha$  whose subjacent functor preserves products. The full subcategory  $\tilde{\mathcal{M}}$  is a variety in  $\mathcal{M}$ .*

*Proof.* (i) Let  $f : (x, \pi_x) \rightarrow (x', \pi')$  be an epimorphism in  $\mathcal{M}$ . Consider the diagram

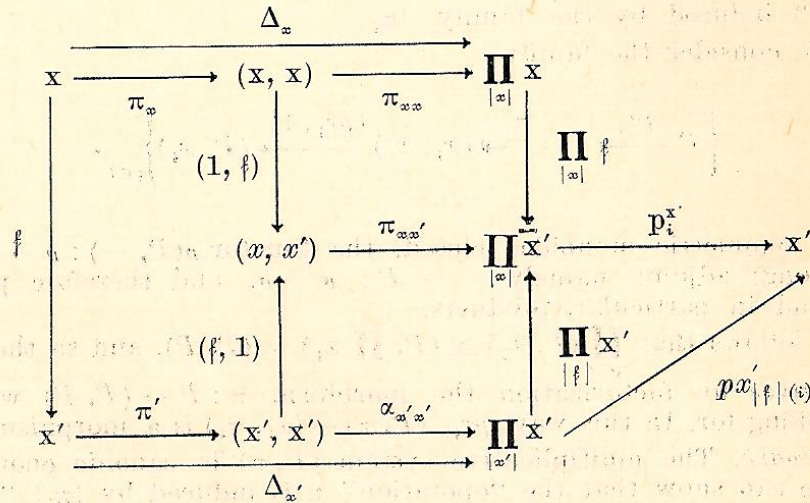


Diagram 4

where  $i \in |x|$ .

Since  $\alpha$  is natural, and  $f$  is a morphism in  $\mathcal{M}$ , the pentagon, the three squares and the triangle are commutative. The two top lines are equal by the definition of  $\pi_x$ . Hence, the bottom lines are also equal since  $f$  is surjective (and so  $\{p_{|f|}^{x'}(i)\}_{i \in |x|}$  are all the projections from  $\prod_{|x'|} x'$ ) and  $f$  is an epimorphism, and we have  $\Delta_{x'} = \alpha_{x'x'} \pi'$ . Thus  $\pi' = \pi_{x'}$  and  $(x', \pi') \in \tilde{\mathcal{M}}$ .

(ii) Now, let  $f : (x, \pi) \rightarrow (x', \pi_{x'})$  be a monomorphism in  $\mathcal{M}$ . Consider the similar diagram

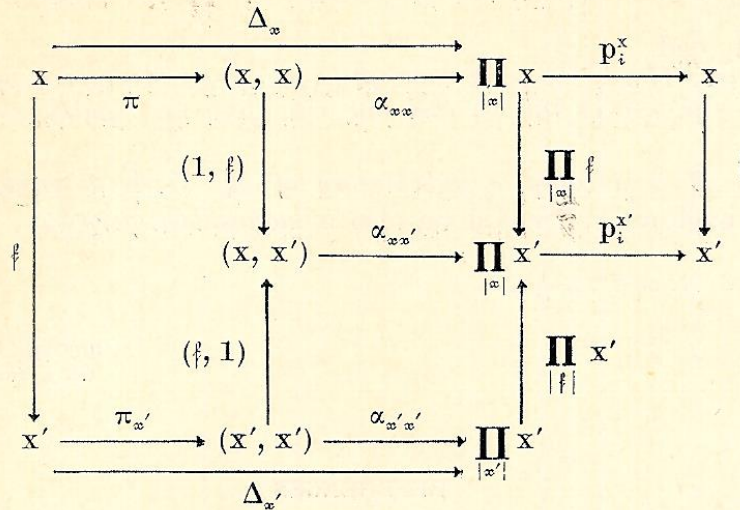


Diagram 5

where  $i \in |x|$ .

The subdiagrams are commutative by similar arguments. The bottom lines are equal by the definition of  $\pi_{x'}$ . Hence the top lines are equal, i.e.  $\Delta_x = \alpha_{xx}\pi$ , since  $f$  is monic and  $\{p_i^x\}_{i \in |x|}$  are all the projections from the product  $\prod_{|x|} x$ . Thus  $\pi = \pi_x$  and  $(x, \pi) \in \tilde{\mathcal{M}}$ .

(iii) We finally consider the diagram

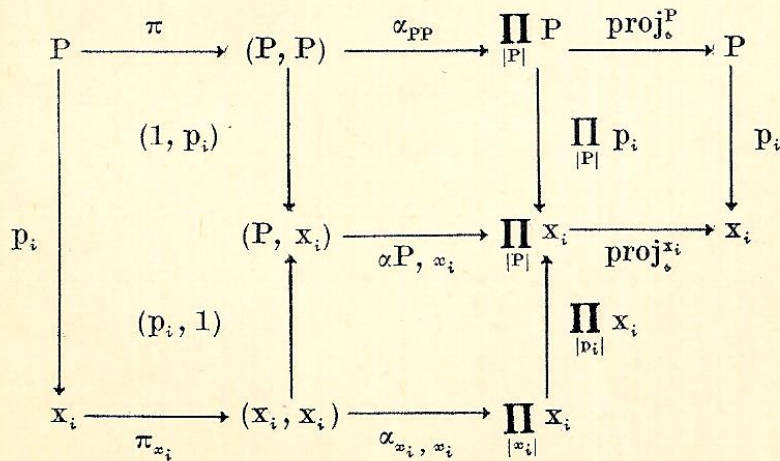


Diagram 6



where  $(x_i, \pi_{x_i}) \in \mathcal{M}$ ,  $P$  has the projections  $p_i$  and  $s \in |P|$ . The top lines are equal and as the subdiagrams are again commutative, the bottom lines are equal, i.e.  $\Delta_P = \alpha_{PP}\pi$ , since  $\{p_i\}_{i \in I}$  are all the projections of  $P$  and  $\{\text{proj}_s^P\}$ ,  $s \in |P|$  are all the projections of  $\prod_{|P|} P$ . Thus  $\pi = \pi_P$  and

$(P, \pi) \in \tilde{\mathcal{M}}$ , qed.

*Remark.* If  $\alpha$  is complete, then  $\mathcal{M}$  is easily seen to be complete because  $\alpha \lim_{i \in I} x_i, -$  preserves limits) and one can give an analogous proof to

show that  $\tilde{\mathcal{M}}$  is a complete subcategory of  $\mathcal{M}$ . Without considering the injective envelopes,  $\tilde{\mathcal{M}}$  could be also a monosubcategory (5).

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